

# Towards a case-based reasoning approach for cloud provisioning

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**Abstract:** Resource provisioning is an important issue of cloud computing. Most of the recent cloud solutions implement a simple way with static thresholds to provide resources. Some more sophisticated approaches consider the cloud provisioning problem a multi-dimensional optimization approach. However, the calculation effort for solving optimization problems is significant. An intelligent resource provisioning with a reduced calculation effort requires smart cloud management methods. In this position paper, we propose a case-based reasoning approach for cloud management. A case records a problem situation in cloud management and its solution. We introduce a case model and a retrieval method for previously solved problem cases with the aim to reuse their re-configuration actions for a recent problem situation. The case model uses the container notion correlated with QoS problems. We present a novel, composite similarity function that allows to compare a recent problem situation with the cases from the past. During retrieval, the similarity function creates a ranking of the cases according to their relevance to the current problem situation. Further, we describe the prototypical implementation of the core elements of our case based-reasoning concept. The plausibility of the retrieval approach has been tested by means of sample cases with simulated data.

## 1 INTRODUCTION

The management of resources for services is a vital aspect of cloud computing (Baun et al., 2011). For a cloud provider, it is important to fulfill the requests of their users and to avoid violations of Service Level Agreements (SLA). On the other hand, it is desirable to use the available resources as effectively as possible to avoid the waste of energy and to save hardware costs. It is required to find a good balance between over-provisioning of resources and under-provisioning (Armbrust et al., 2010). *Cloud management* addresses monitoring and configuration methods to achieve a good cloud configuration. A *cloud configuration* is the placement of the virtual machines (VM) on their physical machines (PM) and the services on the VM's. It also includes the resource provisioning for the VM's and their containers. A *container* is a run time environment such as OpenShift Gears or a Docker container. Cloud management can also be seen as a resource management problem that can be solved by a multi-dimensional optimization approach (Garg et al., 2011; Zhao and Li, 2013), balancing resource provisioning with other optimization criteria like costs for SLA violations. A simple example for over-provisioning is two PM's,

each hosting VM's and with an overall average resource utilization of 30%. It could be beneficial to migrate the VM from one PM to another for the reason of energy savings by stopping one PM. On the other hand, under-provisioning of resources such as memory or bandwidth can lead to SLA violations. Such violations increase the costs for the provider and may lead to a decrease of reputation (Shoaib and Das, 2014).

One method to provide resources is the static way. In this case, the system does not adjust itself to a changing situation. There is the risk of under- or over-provisioning (Shoaib and Das, 2014). Several approaches solve the provisioning problem in a more dynamic way. They range from rather simple, rule based approaches, such as observations on the number of open connections (Pousty and Miller, 2014), to complex algorithms. Quiroz et al. (Quiroz et al., 2009) describe a decentralized, robust online clustering approach for a dynamic set of heterogeneous applications for resource provisioning. They consider the typical workload of a cloud infrastructure as a mix of long running, computationally intensive jobs, bursty and response-time sensitive web service requests, and data and IO-intensive analytic tasks.

*Case-based Reasoning (CBR)* (Aamodt and Plaza, 1994) is an intelligent alternative to the static and dynamic approaches mentioned above. The core idea is to retrieve similar situations and their solutions from the past in order to reuse them for the current situation. CBR has been considered for intelligent cloud management recently in the literature (Maurer et al., 2013). The work of Maurer et al. applies CBR to implement automatic cloud management following the MAPE reference model (Monitor - Analyse - Plan - Execute) (Corporation, 2006), which originates in autonomic computing. A *case* in cloud management records a cloud configuration with current services and SLA's to be processed as a problem situation. A solution describes the optimal distribution of work on the optimal number and configuration of cloud resources while maintaining SLA's. Maurer et al. use a bag of workloads to schedule the work, which makes it difficult to predict future workloads and system behavior.

We think a CBR based approach can solve the provisioning problem in a intelligent way with reduced calculation effort. We are planning to implement a system that uses CBR. This position paper focuses on a core concept for intelligent, case-based cloud provisioning namely the case-based retrieval of former cloud configurations. In our future work, successful cloud configurations from the past shall be reused for problem situations in current cloud configurations.

In CBR, the retrieval "*deals with the process of selecting experience items from the experience base that are relevant for the current problem to be solved.*"(Bergmann, 2002). In this paper, the experience items (cases) are containers with their parent node (the VM that executes the container) and their related services and SLA's. The experience base (case base) contains items with violated SLA's. At this point the cases only include configurations with under-provisioning.

## 2 SERVICE CHARACTERIZATION

In this section we introduce the characterization of services, which is a prerequisite to determine the similarity between two cases. Inspired by Quiroz et al. (Quiroz et al., 2009), we use five attributes to characterize a service. These attributes describe the behavior of the service as follows.

As the name indicates for a *response-time sensitive* service, the response time is important. A

|   | service characterization                |
|---|---|
| 1 | response-time sensitive                 |
| 2 | long running, computationally intensive |
| 3 | bursty, computationally intensive       |
| 4 | data intensive                          |
| 5 | I/O-intensive                           |

Table 1: Overview of service characterization following Quiroz et al. (Quiroz et al., 2009)

scenario for this characterization could be a task with a user interaction where a long waiting time is not feasible.

A *long running, computationally intensive* service is CPU and/or memory intensive and has a long execution time, for example more than 1 hour. In contrast, a *bursty, computationally intensive* service is CPU and/or memory intensive but has a short execution time. A *storage intensive* service needs a large amount of disk space. This could be for example a service that processes large video files and stores them. Such a service would be also *I/O-intensive*. This means the service has special requirements for network bandwidth and/or disk read and write speed.

We have chosen this approach for describing a service instead of the more common *cpu intensive, memory intensive, network intensive and storage intensive*. The reason is, we think this is a more natural way to characterize a service and it describes in a better way what the behavior of a service is. For example the term "the service is cpu intensive" contains less information about a service than the term "the service is a long running, computationally intensive service". On the one hand, cpu intensive does not provide information on the execution time. On the other hand, it could be difficult to determine the exact resource usage. Thus, a more fuzzy description could be suitable.

We are aware of the fact that not all of the service characterizations are in any case independent from each other, and we will consider this in our similarity model. However, we believe that the traditional model can also suffer from similar dependencies. For example, a CPU intensive job will also be memory intensive and /or storage intensive in terms of read/write operations. This may be caused by the fact that the cache is not large enough and the CPU stores intermediate results in the memory or on the disk. This is clearly a dependency from the CPU intensive aspect to the memory intensive aspect.

The similarity function might use alternative service characterizations. We have chosen the above attributes for a service characterization as a starting point for our experiments.

We use a binary vector to determine the characterization of each service. This is the *characterization vector*. For example, a web service that automatically renders large images and stores them has the characterization of *long running*, *computationally intensive*, *storage intensive* and *I/O-intensive*. There is no user-interaction, i.e. the service is not *response-time sensitive*. Because of the large image sizes to be expected it will likely not be *bursty*, *computationally intensive*. The characterization vector for this service is (0,1,0,1,1) following the order described in Table 1. The values of the characterization vector are determined automatically for each service. Initially, we derive default values from the SLA specification (compare Table 2). After a period of monitoring, the characterization of the services may be updated based on observations of the run-time behavior.

| Service characterization                | Related SLA's                         |
|---|---------------------------------------|
| response-time sensitive                 | network latency, bandwidth            |
| long running, computationally intensive | completion time with a long deadline  |
| bursty, computationally intensive       | completion time with a short deadline |
| data intensive                          | data availability                     |
| I/O-intensive                           | network bandwidth, disk IO bandwidth  |

Table 2: Overview of service characterizations and their related SLA's

Garg et al. (Garg et al., 2011) says that transactional applications such as Web applications require guaranteed response time and throughput. We think that related SLA's can determine the network latency and network bandwidth because both values have a strong impact on the response time and the throughput. If the network latency is too high, the response time may also be too high. On the other hand, if the bandwidth is too low the throughput may be affected. The values network latency and bandwidth can be measured by network monitoring tools. The other SLA's are inspired by (Kolodner et al., 2011; Kundu et al., 2010).

### 3 CBR CASE STRUCTURE

First of all we want to give a short introduction into CBR cases before we introduce the separate parts of our model. Afterward we describe how we define the similarity between cases.

A case consists of a problem part and a solution part. The problem part describes a situation that occurred in the past, for example, a cloud configuration with violated SLA's. The solution part contains reconfiguration steps (such as start more VM's, migrate containers etc.) to solve the problem. If a new situation

with violated SLA's occurs, the system will search in the case base to retrieve a case from the past that is similar to the current situation. We use the similarity function we will introduce later on in this section to determine the similarity between two cases. If a similar case is found the solution from the past is used as a starting point for reconfiguration. The problem part of the case includes a container with its parent node, the services executed on the node, the SLA's related to the services, and the service query. The service query is the set of services to be started next. They have not yet been provided with resources.

There are two possible problems that can occur, which may make it necessary to reconfigure the cloud configuration. The first problem is that there is a container which contains one or more violated SLA's, or SLA's that are about to be violated in the near future. The second problem is that there are service requests that can not be fulfilled by the node because of its resource utilization. In both cases, a reconfiguration may be required.

We describe a *cloud configuration C* as a set of nodes  $n \in Nodes$ . A node can be a virtual machine or a storage component. The nodes form the first layer of a hierarchical structure as depicted in Figure 1. The second layer comprises the containers and the third the services.

Each node  $n \in Nodes$  includes the provided resources  $pr^{node}$ , the current utilization  $ur^{node}$  of each resource and a set of containers that are executed by this node  $n = (pr^{node}, ur^{node}, containers)$ . The provided resources depend on the type of node. The provided resources of a virtual machine are described by values for CPU, memory, storage and network bandwidth  $pr^{node} = (cpu, mem, sto, bdw)$ . The resources can be monitored by different tools, such as nagios (Nagios, ). The utilization of the resources is given as the resource usage (in percentage) of the containers that are executed on this node. All containers at the same node share the resources of this node. Thus, the resource utilization of the node can be determined from the resource utilization of the particular containers.

Similar to nodes, containers are run time environments for services. They are described by the provided resources  $pr^{con}$ , utilized resources  $ur^{con}$ , and a set of services, which use the provided resources and are responsible of resource utilization in analogy to containers for nodes. Let  $c$  be a container of node  $n$ , then is  $c = (pr^{con}, ur^{con}, services)$ .

Each service can be considered a *workload*, consisting of a unique service identifier, the service characterization, the number of users who are currently working with, and the volume of the input data. An example of a service is a rendering Web service where

users can upload their images and let them being rendered. The service characterization of such a service will probably be *response-time sensitive* due to the user interaction, and *bursty, computationally intensive* respectively *long running, computationally intensive* due to the broad diversity of image sizes. Such a service could be used by several users at the same time. The number of the users who currently communicate with the service can be discovered by the different IP addresses. The uploaded images are the input data and the size (in mb) is the volume of the input data. This can be monitored by network-traffic monitoring tools.

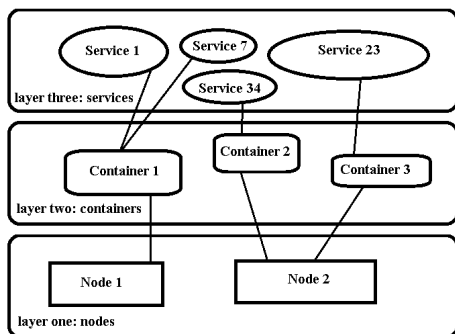


Figure 1: The hierarchic structure of cloud configuration.

One challenge is to assign the resource usage of a node to it's containers and likewise from a container to it's services. For example, if the memory utilization of a node  $n$  is 66% and  $n$  comprises only one container  $c_1$  it is easy to determine the share of the memory utilization consumed by  $c_1$ . Let us assume that  $n$  provides  $c_1$  with all available resources, i.e.  $pr^{node} = pr^{con}$ . In this case the 66% memory utilization comes all from  $c_1$ . If  $n$  comprises three containers ( $c_1, c_2, c_3$ ) the case is more difficult. Even if the resources  $pr^{node}$  of  $n$  are distributed in equal parts between the containers (for each  $c_i, pr^{con} = 1/3 pr^{node}$ ) it can not be concluded that the containers have the same workloads and, thus, are responsible for 22% of the overall utilization. It is possible that one of the containers is idle while another one has a share of 11% and the last container is responsible for 55% of the utilization as illustrated in Figure 2. To estimate the percentage of utilization consumed by each container /service, we use task manager tools. At each level of the hierarchy (node, container, service), the current share of the CPU, memory, network and disk I/O usage of the containers/services can be determined via task manager tools like top (Unix Top, 2014) which show the individual share of the resources for each process. Each container is visible as a process for a node and each service for a container.

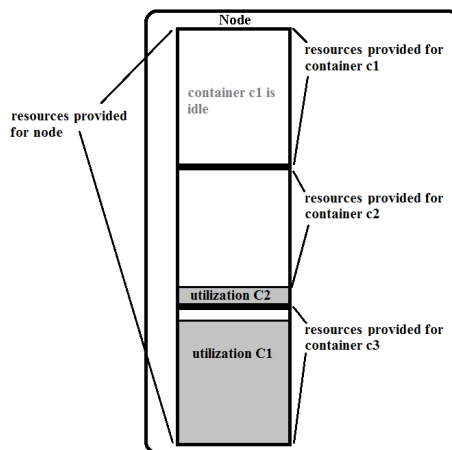


Figure 2: Example of resources provided and utilized.

In addition to the cloud configuration, our model of a problem situation considers the *service query*. Each time a cloud user considers to start a service she sends a request to the cloud system to trigger this service. The *service query* contains all requests that have not yet been fulfilled. The requests are described by the name of the service and the according SLA's. A sample service is the rendering Web service *Render Service on Apache2.4*. After a request is accomplished, the system assigns a unique identifier to the service and deletes the request from the *service query*. The status of the individual SLA's for ongoing services are monitored by the system. The state of a SLA might be green, yellow or red. Green means the measured values are acceptable im terms of the SLA. If the state is yellow the values are still acceptable but they are close to boarder and a SLA violation in the near future is to be expected. A red state means that an SLA violation has occurred. The thresholds forchangig the state from green to yellow or to red depend on the SLA, i.e. it is predefined by the administrator or the user.

## 4 CASE SIMILARITY

In order to determine the similarity between two cases, we use a composite similarity function for the problem part of the cases. We define the similarity between two nodes by the similarity of their resources provided, including the utilization of these resources. The similarity function for nodes  $nodeprov$  is induced from a taxonomy of nodes as depicted in Figure 3. For our first experiments, we use the resource sets of Amazon EC2 instances (AWS, ) as leafs of the taxonomy. An example for an EC2 instance is the

M3 instance with two virtual CPU's, 3.75 GB RAM and 32 GB of disk space. For further information on the instances, we refer to the Amazon Web page [www.aws.amazon.com/de/ec2/instance-types/](http://www.aws.amazon.com/de/ec2/instance-types/). We think that the size of the instances (tiny, small, 4xlarge and so on) has a higher impact on the similarity than an eventual specialization of the container (for example for memory intensive applications). For the nodes, we prefer a taxonomical similarity over a numerical function calculated from the Euclidean distance, for instance. The reason is that for many nodes the values between several resources differ to a large extent. For example, the minimal number of CPU's is one and the maximum number is 40. On the other hand, the minimal bandwidth we expect is 100 Mbits/sec. The maximum is 10 Gbits/sec. Instead of balancing the similarity values for the particular resources by weighting or by normalization, we decided to build a taxonomy since it provides an easy and natural model. The similarity of the resource utilization *nodeutil* is

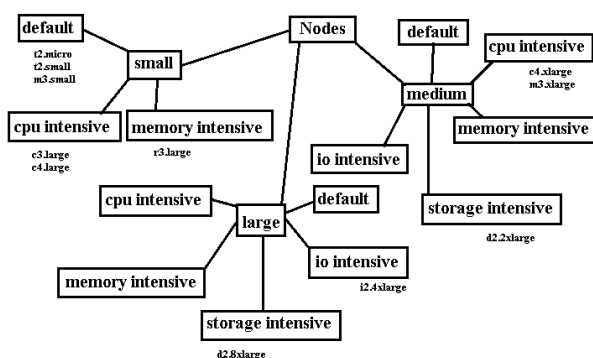


Figure 3: Taxonomy of nodes (in parts).

calculated by the Euclidean distance  $nodeutil(p, q) = \sqrt{\sum_{i=1}^n (q_i^{node} - p_i^{node})^2}$  where  $p$  is the vector of  $n$  utilization values for the first case and  $q$  for the second case. The utilization values are provided in percentage. For example,  $q_1^{node} = 50$  is the utilization of the CPU  $q_1^{node}$  with a value of 50%.  $p_2$  is the utilization of the memory and so on.

Analog to nodes, we specify the similarity function between containers *consim* by means of a taxonomy of the provided resources and the Euclidean distance for the utilized resources *conuti*. We use RedHat OpenShift Gears and IBM Bluemix containers as templates for our containers. Thus, the similarity function for containers is  $consim(p, q) = \sqrt{\sum_{i=1}^n (q_i^{con} - p_i^{con})^2}$  with *con* for container.

In addition, we consider the services and the SLA's related to the container. As pointed out earlier, every service has a vector of it's characterizations (see Section 2). To compare two nodes it is not necessary to compare the services themselves but to compare the occurrence of a characterization attribute, i.e. whether a service is response-time sensitive or not. Therefore, we build the container characterization *conchar*. This is a vector of aggregated occurrences of each service characterization attribute for each service that is executed on the container. For example: a container  $c$  has two services  $s1, s2$ . Both services are *response-time sensitive*. In addition,  $s2$  is also *bursty*, *computationally intensive*. The container characterization vector  $CV$  for  $c$  is then  $CV = (2, 0, 1, 0, 0)$ . We use again the Euclidean distance for the container characterization vectors.

To determine the similarity of the number of SLA violations *slavio* we build again a vector that contains every SLA type occurring in our cloud configuration and sum up the number of SLA's with condition red for the specific container. For example, there might be two different types of SLA's specified in our cloud configuration namely network latency and data availability. Again, we consider container  $c$  with two services  $s1, s2$ . Both services contain an SLA on network latency and data availability. In this example, both SLA's are violated for  $s1$ . Due to different thresholds, only the network latency SLA for  $s2$  is in red state. Thus, the resulting SLA violation vector for  $c$  is  $c = (2, 1)$ . Again, we use the Euclidean distance to measure the similarity between two SLA violation vectors.

Finally, we compare the service query  $sq$  for both cases. At the moment, we compare only the length of two queries to determine the similarity. In future work, we will extend this.

Now we can calculate the entire similarity *osim* for two cases by aggregating the local similarity values in a weighted sum:

$$osim = x_1 * nodesim + x_2 * nodeuti + x_3 * consim + x_4 * conuti + x_5 * conchar + x_6 * slavio + x_7 * sq.$$

The weights  $x_1, x_2, \dots, x_7$  are configurable.

## 5 EVALUATION

In a preliminary evaluation, we have implemented a prototype to generate 50 test samples and measured the similarity between them. As mentioned before, we have used Amazon EC2 instances as a template for our nodes and OpenShift Gears and Bluemix containers as a template for our containers. We used about 100 automatically generated services with a random

set of service characterizations. The nodes for the cases have been chosen randomly with random sets of containers and random sets of services. Depending on the service characterizations, we determined randomly a utilization value for each resource (CPU, memory...). For example, if a service is long running, computationally intensive, the average CPU and memory utilization will be higher than for a service without this characterization. The utilization is also higher for a container with multiple characterizations of the same type. For example, if a container contains two services with the long running, computationally intensive characterization the entire utilization value will be higher than for a single one. After having initialized the utilization for each container, we calculated the resource utilization for the node. Further, a set of SLA's is assigned to the containers, depending on the services executed on the container. If an SLA can not be fulfilled due to the resource utilization of the container, the SLA is set to status red. The case base has been created by choosing randomly one container per node.

We applied the similarity function described in Section 4. The results in Figure 4 indicate that the size of the nodes (for example Case 4 is an i2.xlarge node) is an important aspect since nodes with a similar size are frequently more similar to each other than two nodes with other sizes. The result shows also that the distance of a case to itself is always zero, i.e. that the cases are equal.

| Case:            | Case1 c3.2xlarge | Case2 c4.xlarge | Case3 r3.4xlarge | Case4 i2.xlarge | Case5 i2.4xlarge |
|------------------|------------------|-----------------|------------------|-----------------|------------------|
| Case1 c3.2xlarge | 0.0              | 6857.0          | 16080.0          | 8188.0          | 21039.0          |
| Case2 c4.xlarge  | 6857.0           | 0.0             | 16496.0          | 9465.0          | 19721.0          |
| Case3 r3.4xlarge | 16080.0          | 16496.0         | 0.0              | 18959.0         | 10630.0          |
| Case4 i2.xlarge  | 8188.0           | 9465.0          | 18959.0          | 0.0             | 25419.0          |
| Case5 i2.4xlarge | 21039.0          | 19721.0         | 10630.0          | 25419.0         | 0.0              |

Figure 4: Example result of our experiments.

## 6 CONCLUSION

In this paper, we have presented our concept for the retrieval part of a case-based approach for intelligent cloud provisioning. We have introduced our similarity function for cases and have conducted several test evaluations. The evaluation with sample test cases has shown that it is possible to retrieve plausible results. There are several open issues we will tackle in future. We have only considered under-provisioning so far. However, we will develop a CBR approach for preventing over-provisioning as well. Second, we will consider whether it is sufficient to observe only single containers with their nodes. Another open issue is to determine the characterization for services based on their run-time behavior. The preliminary results are

promising and provide a first, important step towards intelligent, case-based cloud provisioning.

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